The Landscape of Computing Symmetric n-Variable Functions with 2n Cards

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#### Introduction

# Secure Multi-Party Computation

- Alice and Bob want to know if they both like each other.
  - No one wants to confess first.
- Needs a protocol that only distinguishes between the two cases: they both like each other, and anything else.



# Secure Multi-Party Computation

- Each having a bit *a* and *b* of either 0 or 1.
- Needs a protocol that computes  $a \wedge b$  without leaking unnecessary information.
  - If a player's bit is 1, he/she inevitably knows other player's bit.
  - If a player's bit is 0, he/she should know nothing about other player's bit.

# **Card-based Protocols**

- Protocols using physical cards
- Does not require computer
- Uses only small, portable objects
- Easy for observers to verify the correctness and security, even for nonexperts
- Suitable for teaching purpose



#### The Five-Card Trick

- Developed by den Boer in 1989, beginning of the study in card-based cryptography.
- Uses five cards: three identical & cards and two identical 
  cards.
- Encodes 0 by ♣♥ and 1 by ♥♣.

#### The Five-Card Trick

- Each player is given one \* and one \*, with another \* (helping card) faced down on the middle of table.
- Alice places her cards encoding *a* to the left of the middle card.
- Bob places his cards encoding b to the right of the middle card.



Then, we swap the fourth and the fifth cards.



Observe the cyclic rotation of the deck.

# The Five-Card Trick

- We obscure the initial position of the cards by shuffling the deck into a random uniform cyclic permutation.
  - i.e. a permutation uniformly chosen from  $\{id, (12345), (12345)^2, (12345)^3, (12345)^4\}$
- Can distinguish the case a = b = 1 from other cases.

# The Four-Card Trick

 Later, in 2012, Mizuki et al. showed that the AND function can be computed with four cards, using no helping card.



#### **Other Functions**

- Besides the AND function, protocols to compute other Boolean functions have also been developed.
- In 2009, Mizuki and Sone developed a fourcard XOR protocol.



#### **Other Functions**

- As each input bit is encoded by two cards, computing an *n*-variable function requires at least 2*n* cards.
- Any *n*-variable symmetric function can be computed with 2n + 2 cards (Nishida et al., 2015).
- We are interested in optimal protocols that use exactly 2n cards.

# **Properties of Protocols**

- Number of shuffles as low as possible
- Committed format output is in the same format as input (♣♥ for 0 and ♥♣ for 1)

• A Boolean function  $f: \{0,1\}^n \rightarrow \{0,1\}^n$  is called symmetric if

 $f(x_1, \dots, x_n) = f(x_{\sigma(1)}, \dots, x_{\sigma(n)})$ for any  $x_1, \dots, x_n$  and any permutation  $(\sigma(1), \dots, \sigma(n))$  of  $(1, \dots, n)$ .

• Note that the output value only depends on the sum  $\sum_{i=1}^{n} x_i$ .

• Denote an n-variable symmetric Boolean function by  $S_X^n$  for some  $X \subseteq \{0, ..., n\}$ .

• 
$$S_X^n = \begin{cases} 1, & \text{if } \sum_{i=1}^n x_i \in X; \\ 0, & \text{otherwise} \end{cases}$$

• E.g.  $x_1 \wedge \cdots \wedge x_n$  is denoted by  $S_{\{n\}}^n$ .

- If f is computable by a number of cards,
  - Negating all variables in f is also computable by the same number of cards.
  - $\circ$  So is the negation of f.
- We can classify all *n*-variable symmetric functions into several NPN-equivalence classes.
- $S_X^n$  is in the same class as  $S_{\{0,\dots,n\}-X}^n$  and  $S_{\{n-x|x\in X\}}^n$ .

# Summary of Protocols



## **Two Variables**

- Eight functions
- Three classes
- Trivial (constant), AND, XOR



#### **Two Variables**

Function	Name	Protocol	Comm itted?	#Shuf- fles	Other Functions in the Same Class
$S_{\phi}^{3}$	Constant	trivial			$S^3_{\{0,1,2\}}$
$S^{3}_{\{1\}}$	XOR	Mizuki-Sone, 2009	>	I	$S^3_{\{0,2\}}$
$S^{3}_{\{2\}}$	AND	Mizuki et al., 2012	X	2	$S^3_{\{0\}}, S^3_{\{0,1\}}, S^3_{\{1,2\}}$

 No four-card committed-format AND protocol with finite number of shuffles (Koch et al., 2015)



# **Three Variables**

- I6 functions
- Six classes (one is trivial)
- Fully solved in 2023.



#### **Three Variables**

Function	Name	Protocol	Comm itted?	#Shuf- fles	Other Functions in the Same Class
$S_{\phi}^{3}$	Constant	trivial	$S^3_{\{0,1,2,3\}}$		
$S^{3}_{\{1,3\}}$	XOR	Mizuki-Sone, 2009	<b>~</b>	2	$S^3_{\{0,2\}}$
$S^{3}_{\{3\}}$	AND	Mizuki, 2016	X	5	$S^3_{\{0\}}, S^3_{\{0,1,2\}}, S^3_{\{1,2,3\}}$
		Isuzugawa et al., 2021	X	2	
$S^{3}_{\{0,3\}}$	Equality	Shinagawa-Mizuki, 2019	X		$S^{3}_{\{1,2\}}$
		R-Itoh, 2021	<b>~</b>	2	
$S^{3}_{\{2,3\}}$	Majority	Toyoda et al., 2021	X	2	$S^3_{\{0,1\}}$
$S^{3}_{\{1\}}$	-	Shikata et al., 2023	X	3	$S^3_{\{2\}}, S^3_{\{0,1,3\}}, S^3_{\{0,2,3\}}$



#### **Four Variables**

- 32 functions
- Ten classes (one is trivial)
- Eight currently have protocols (two of them are Las Vegas protocols)
- Two open problems



#### **Four Variables**

Function	Name	Protocol	Comm itted?	#Shuf- fles	Other Functions in the Same Class
$S_{\phi}^{3}$	Constant	trivial			$S^3_{\{0,1,2,3\}}$
$S^{3}_{\{1,3\}}$	XOR	Mizuki-Sone, 2009	<b>~</b>	3	$S^3_{\{0,2,4\}}$
$S^{3}_{\{4\}}$	AND	Mizuki, 2016	X	5	$S^3_{\{0\}}, S^3_{\{0,1,2,3\}}, S^3_{\{1,2,3,4\}}$
$S^{3}_{\{0,4\}}$	-	R-Itoh, 202 I	<b>~</b>	3	$S^3_{\{1,2,3\}}$
$S^{3}_{\{2\}}$	-		X	3	$S^3_{\{0,1,3,4\}}$
$S^{3}_{\{1\}}$	-	Shikata et al., 2023	X	≈7	$S^3_{\{3\}}, S^3_{\{0,1,2,4\}}, S^3_{\{0,2,3,4\}}$
$S^{3}_{\{1,2\}}$	-		X	≈8	$S^3_{\{2,3\}}, S^3_{\{0,1,4\}}, S^3_{\{0,2,4\}}$
$S^{3}_{\{0,3\}}$	Div3	R, 2023	X	4	$S^3_{\{1,4\}}, S^3_{\{0,2,3\}}, S^3_{\{1,2,4\}}$
$S^{3}_{\{3,4\}}$	Majority	Open problem			$S^3_{\{0,1\}}, S^3_{\{0,1,2\}}, S^3_{\{2,3,4\}}$
$S^{3}_{\{0,2\}}$	-				$S^3_{\{2,4\}}, S^3_{\{0,1,3\}}, S^3_{\{1,3,4\}}$

# More than Four Variables

- In 2022, Shikata et al. proved that there exists a 2n-card protocol to compute any n-variable symmetric function with  $n \ge 8$ .
- Limits the open problems to n = 4,5,6,7.

# More than Four Variables

- *n* = 5
  - 64 functions, 20 classes
  - 7 solved, 13 open
- *n* = 6
  - 128 functions, 36 classes
  - 10 solved, 26 open
- *n* = 7
  - 256 functions, 72 classes
  - 14 solved, 58 open

# More than Four Variables

- The number of NPN-equivalence classes is the number of n-bead black-white reversible strings.
- Follows the sequence A005418 in OEIS.
- 1, 2, 3, 6, 10, 20, 36, 72, 136, 272, 528, 1056, ...

#### **Questions and Comments**