

TRANSLATION OF THEOREM-PROVING TEXT IN $MTSR$ TO NATURAL
LANGUAGE TEXT

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Abstract. In the article is described translation of theorem-proving text from $MTSR$ language into natural language. As example, it is considered translation into English. Used method is valid for any language.

Keywords and phrases: $MTSR$ theory, theorem-proving text, formal grammar, Bison program, parsing.

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1. Introduction. In the article is described translation of theorem-proving text from $MTSR$ language into natural language. As example, it is considered translation into English. Used method is valid for any language. Source language's grammar is traduced into LR (1) grammar [1] and by the grammar is composed Bison program. Compilation of Bison program gives working program. Input of the program is theorem-proving text and output is its transfer into target language's text [2-3].

2. TSR -Logic Language. The language of TSR logic consists of the following symbols:

1. Fundamental symbols:

(a) Logical connectives: \neg (of the weight 1), \wedge , \vee , \rightarrow , \leftrightarrow (each of the weight 2).

(b) Logical operational sign \mathcal{T} of the weight (1, 1).

(c) Substantive special substitution operator S of the weight (1, 2).

(d) Relational logico-special substitution operator R of the weight (1, 2) and with the logicity indicator 2.

(e) Object letters: X_0, X_1, \dots (1)

(f) Predicate symbols $=$ and \in , each of the weight 2, and predicate letters: $P_0^n, Q_0^n, P_1^n, Q_1^n, \dots$ (2)

(g) Functional symbol \supset that has the weight 2, and functional letters: $f_0^n, g_0^n, f_1^n, g_1^n, \dots$ (3)

(h) [and] (left and right brackets)

2. Signs, introduced by the definitions of the types I, II and II'.

Finite sequence of signs of TSR are called a *word* of TSR logic. The words $\mathcal{T}X_0, \mathcal{T}X_1, \dots$ are the TSR logic operators with the weight 1. The words SX_0, SX_1, \dots and RX_0, RX_1, \dots are the TSR logic operators with the weight 2. Besides, the operators SX_0, SX_1, \dots are substantive partial quantifiers with the binding indicator 2, and

the operators RX_0, RX_1, \dots are logico-special partial quantifiers with logicity and binding indicator 2.

The operator is called :

1) A logical relational if its operands are formulas, and results is formulae. For examples: $\neg, \bigvee, \forall x, \exists x \dots$

2) A logical substantive if its operands are formulas, and results is term. For examples: $\mathcal{T}x, \varepsilon x, \dots$

3) A special relational if its operands are terms, and results is formulae. For example: $\langle \text{represent} \rangle x, \dots$

4) A special substantive if its operands are terms and results is term. For examples: $+, \sin, Sx, \dots$

5) A logico-special relational, if its operands are terms and results is formulae. For examples: $Rx, \langle \text{root} \rangle x, \dots$

6) A logico-special substantive if its operands are terms and results is term. For examples: $\langle \text{subset} \rangle x, \dots$

Formulas and *terms* of \mathcal{TSR} logic are defined in following way:

1. Subject letters are simple terms.
2. If σ is an n -ary relational logical (resp. special) operator, then $\sigma A_1 \dots A_n$ (resp. $\sigma T_1 \dots T_n$) is either a formula or a term depending whether σ is relational or substantive.
3. Let C_1, \dots, C_n be a sequence of formulas. If σ is an n -ary logico-special operator whose logicity indicator is (n_1, \dots, n_k) , and C_{n_1}, \dots, C_{n_k} is the maximal subsequence of the sequence C_1, \dots, C_n consisting of formulas only, then $\sigma C_1, \dots, C_n$ is either a formula or a term depending whether σ is relational or substantive.
4. C is formula or term if and only if it is derived by the three rules above.

3. Translation from \mathcal{TSR} Grammar into $LR(1)$ Grammar. Suppose $G = (N, \Sigma, P, S)$, where Σ is set of terminal symbols from \mathcal{TSR} , N is set of non-terminal symbols, S is starting symbol from N , P is set of rules corresponding to rules from \mathcal{TSR} grammar:

$\langle \text{theorem-proving-text} \rangle ::= \langle \text{theorem-statement} \rangle \langle \text{theorem-proving-with-comments} \rangle$
 $\langle \text{theorem-statement} \rangle ::= \langle \text{label} \rangle M\mathcal{TSR} \langle \text{formula} \rangle$
 $| \langle \text{label} \rangle M\mathcal{TSR} \langle \text{sequence-formula} \rangle \langle \text{formula} \rangle | \langle \text{label} \rangle (\langle \text{formula1} \rangle$
 $\langle \text{complement-symbol} \rangle \langle \text{formula1} \rangle \langle \text{complement-symbol} \rangle M\mathcal{TSR} \langle \text{formula} \rangle$
 $\langle \text{formula1} \rangle ::= M\mathcal{TSR} | \langle \text{sequence-formula} \rangle | \langle \text{formula} \rangle$
 $\langle \text{label} \rangle ::= ((C|ID|I) \langle \text{digits} \rangle | (ra|rb|rc|rd|re)).$
 $\langle \text{digits1} \rangle ::= \langle \text{digit1} \rangle \langle \text{digits} \rangle | \langle \text{digit1} \rangle$
 $\langle \text{digits} \rangle ::= \langle \text{digit1} \rangle | 0$
 $\langle \text{digit1} \rangle ::= 1|2|3|4|5|6|7|8|9$
 $\langle \text{sequence-formulae} \rangle ::= \langle \text{formula} \rangle, \langle \text{sequence-formulae} \rangle | \langle \text{formula} \rangle$
 $\langle \text{formula} \rangle ::= \langle \text{one-place-operator} \rangle \langle \text{formula} \rangle |$
 $(\langle \text{formula} \rangle \langle \text{two-place-operator} \rangle \langle \text{formula} \rangle) |$
 $\langle \text{atom} \rangle$

$\langle \text{one-place-operator} \rangle ::= \emptyset$
 $\langle \text{two-place-operator} \rangle ::= \acute{U} | \acute{U} | \textcircled{R} | \ll$
 $\langle \text{atom} \rangle ::= A \langle \text{digits} \rangle | B \langle \text{digits} \rangle | A | B$
 $\langle \text{theorem-proving-with-comments} \rangle ::= \vdash MTSR, \langle \text{sequence-members} \rangle \dashv MTSR$
 $\langle \text{sequence-members} \rangle ::= \vdash MTSR \text{---} \langle \text{sequence-formulae} \rangle \text{---}, \langle \text{sequence-members} \rangle,$
 $\dashv MTSR |$
 $\langle \text{sequence-formula} \rangle \text{---} \langle \text{formula} \rangle \dots, \langle \text{sequence-members} \rangle |$
 $\langle \text{formula} \rangle.$
 $\langle \text{label} \rangle, \langle \text{sequence-members} \rangle | \langle \text{formula} \rangle .. | \langle \text{formula} \rangle. \langle \text{label} \rangle$
 The grammar is $LR(1)$ grammar.

4. Translation from $LR(1)$ Grammar into BISON Program. Grammar' rules from paragraph 3 we should add semantic rules to receive Bison's program file. Bison program has the following structure:

```

%{
    Prolog
} %
Bison Declaration
%%
input : /* Empty line */
input line
;
line: '\n'
| theorem_proving_text '\n' { printf ("\t %.10g\n", S1); }
;
theorem_proving_text : theorem_statement theorem_proving_with_comments
{ SS = S1 "." S2 ; } ;
theorem_statement : label MTSR formula { SS = S1 "." " " " " " " " " " " " " " " MTSR
" " " " " " " " " " S3 ; }
| label MTSR ( sequence_formula ) formula label ( formula1 complement_symbol
formula1 ) complement_symbol MTSR formula { SS = S1 "." " " MTSR " " " " { " " "
S4 "." " " } " " S6 "." S7 "." " { " " S9 "." S10 "." S11 "." " " } " " S13 "." " MTSR
" " " S15 ; }
formula1: MTSR | sequence_formula | formula
| label: ( (C | ID | I) digits |(ra |rb |rc |rd | re )).
digits1: digit1 digits | digit1
digits: digit1 | 0
digit1: 1|2|3|4|5|6|7|8|9
sequence_formulae: formula, sequence_formulae | formula
formula: one_place_operator formula |
(formula two_place_operator formula )|
atom
one_place_operator : \emptyset
two_place_operator: \acute{U} | \acute{U} | \textcircled{R} | \ll
atom: A digits | B digits | A | B

```

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theorem_proving_with_ comments:  $\vdash MTSR$ , <sequence_ members  $\dashv MTSR$ 
sequence_ members:  $\vdash MTSR$  | sequence_ formulae |, sequence_ members,
 $\dashv MTSR$  |
sequence_ formula || formula .., sequence_ members | formula. label,
sequence_ members | formula ..| formula. label
%%
Epilogue

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5. Getting Target Text. After translation of Bison program we receive working program. Running the program with input (theorem-proving text in TSR) will give corresponding text in target language.

R E F E R E N C E S

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